

Active Learning of Mealy Machines with Timers

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Many real-world scenarios have **timing** constraints, e.g., giving a Highlights talk, network protocols.

¹Angluin, “Learning Regular Sets from Queries and Counterexamples”, 1987.

²Waga, “Active Learning of Deterministic Timed Automata with Myhill-Nerode Style Characterization”, 2023.

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Well-known model for such systems: **timed Mealy machines** which are Mealy machines (*i.e.*, defining a function from input words to output words) augmented with **clocks**.

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Learning (in Angluin's¹ style) timed Mealy machines² is hard.

We focus on **timers** \rightsquigarrow **Mealy machines with timers (MMT)**.

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Clocks start at zero and grow towards infinity.

Timers start a given value and decrease towards zero.

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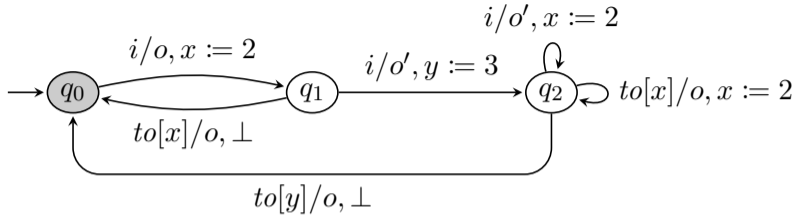


Figure 1: A Mealy machine with timers (**MMT**).

(q_0, \emptyset)

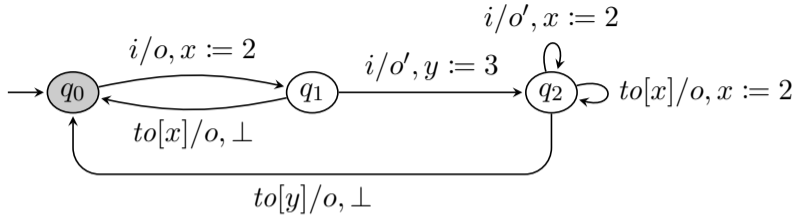


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$$(q_0, \emptyset) \xrightarrow{0.5} (q_0, \emptyset)$$

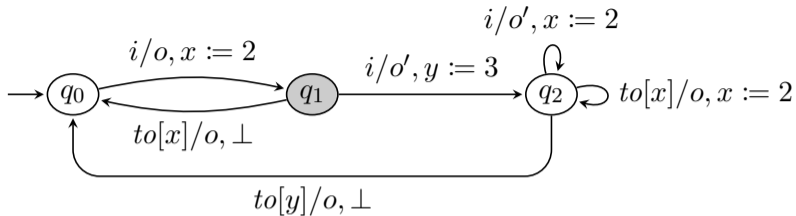


Figure 1: A Mealy machine with timers (**MMT**).

$$(q_0, \emptyset) \xrightarrow{0.5} (q_0, \emptyset) \xrightarrow{i/o} (q_1, x = 2)$$

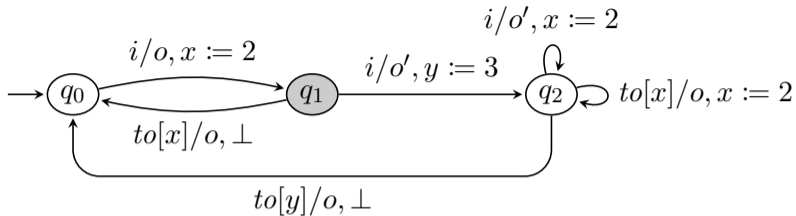


Figure 1: A Mealy machine with timers (**MMT**).

$$(q_0, \emptyset) \xrightarrow{0.5} (q_0, \emptyset) \xrightarrow{i/o} (q_1, x = 2) \xrightarrow{2} (q_1, x = 0)$$

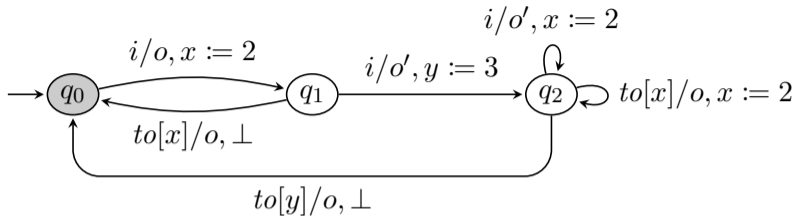


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$$(q_0, \emptyset) \xrightarrow{0.5} (q_0, \emptyset) \xrightarrow{i/o} (q_1, x = 2) \xrightarrow{2} (q_1, x = 0) \xrightarrow{to[x]/o} (q_0, \emptyset)$$

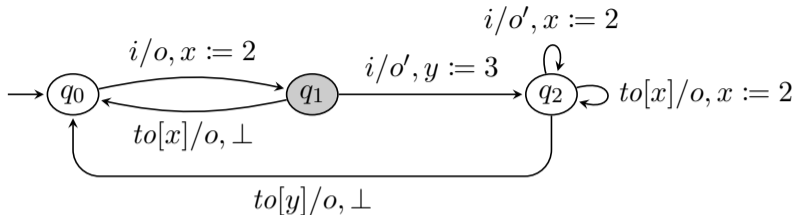


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$$\begin{aligned}
 (q_0, \emptyset) &\xrightarrow{0.5} (q_0, \emptyset) \xrightarrow{i/o} (q_1, x = 2) \xrightarrow{2} (q_1, x = 0) \xrightarrow{to[x]/o} (q_0, \emptyset) \\
 &\xrightarrow{0} (q_0, \emptyset) \xrightarrow{i/o} (q_1, x = 2)
 \end{aligned}$$

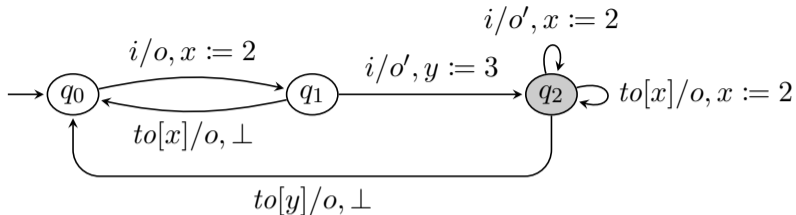


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 &\xrightarrow{0} (q_0, \emptyset) \xrightarrow{i/o} (q_1, x = 2) \xrightarrow{2} (q_1, x = 0) \xrightarrow{i/o'} (q_2, x = 0, y = 3)
 \end{aligned}$$

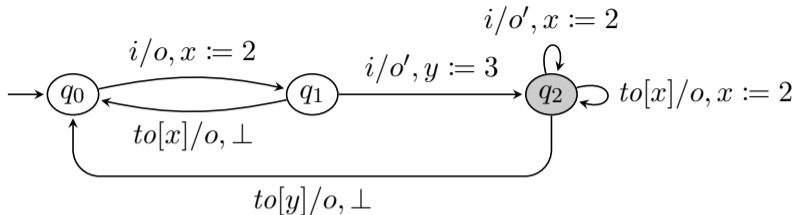


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 &\xrightarrow{0} (q_2, x = 0, y = 3).
 \end{aligned}$$

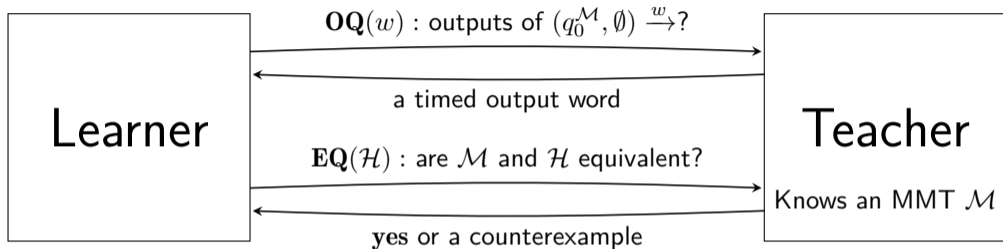


Figure 2: Adaptation of Angluin's framework³ to MMTs.

³Angluin, "Learning Regular Sets from Queries and Counterexamples", 1987.

We adapt $L^\#$ (active learning algorithm for Mealy machines⁴) to MMTs: $L^\#_{\text{MMT}}$.

Theorem 1. Let \mathcal{M} be a **good** MMT and ℓ be the length of the longest counterexample. Our $L^\#_{\text{MMT}}$ learning algorithm requires a number of queries **polynomial** in the number of states of \mathcal{M} and in ℓ , and **exponential** in the number of timers of \mathcal{M} .

⁴Vaandrager et al., “A New Approach for Active Automata Learning Based on Apartness”, 2022.

⁵Bruyère, Pérez, et al., “Automata with Timers”, 2023.

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Problem. There may be actions that occur **simultaneously**, due to **timeouts**.

↔ A **good** MMT ensures that we can observe every behavior by changing some delays.⁵

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⁵Bruyère, Pérez, et al., “Automata with Timers”, 2023.

Thank you!




For all details, see

Bruyère, Pérez, et al., “Automata with Timers”, 2023, and
Bruyère, Garhewal, et al., “Active Learning of Mealy Machines with Timers”, 2024.



For our implementation, see

<https://doi.org/10.5281/zenodo.10647628>, and
https://gitlab.science.ru.nl/bharat/mmt_lsharp.

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